

# Problem Set #5

## IE 495

Due October 30, 2000

### Written Problems

1. Prove that in a binomial heap, the amortized running time for an insert operation is  $O(1)$ . There is an argument for this in the reading by Kozen and we also discussed it in class. You should give a more detailed proof that shows you understand the concept of amortized analysis. The proof approach we discussed in class is a little more rigorous and straightforward than the one in Kozen. You might try to use this approach.
2. Another way of solving the disjoint set union problem is as follows:

Let  $NAME(i)$  be the name of the set containing  $i$ ,  $NUMBER(j)$  the number of items in set  $j$ , and  $LIST(j)$  a pointer to a linked list of the items in set  $j$ . The  $FIND(i)$  operation can then be accomplished simply by examining  $NAME(i)$ . The  $UNION(j, k)$  operation where  $j$  and  $k$  denote sets, is accomplished by first comparing  $NUMBER(j)$  with  $NUMBER(k)$ . If  $NUMBER(j) \leq NUMBER(k)$ , then

- Set  $NAME(i) = k$  for all  $i$  in  $LINK(j)$ ,
- Append  $LINK(j)$  to  $LINK(k)$
- Increase  $NUMBER(k)$  by  $NUMBER(j)$

The new set is named  $j$ . Prove that for a total of  $n$  items, the time for all  $UNION$  operations is  $O(n \log n)$ .

3. Write an algorithm to delete an identifier from a hash table in which overflows are handled by linear probing (see Horowitz and Sahni, page 87).

### Programming Problems

4. This is a group problem. Each project group should only produce one code. Modify your code from the last assignment so that it performs component labeling. In other words, it should read in a graph from a file as a list of edges and print out the list of nodes in each connected component of the graph.