

Sparse Simplex and LPs with Embedded Network Structure

CORAL Seminar Series

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Reference

Gülpinar, N., G. Mitra, and I. Maros 2002. Creating Advanced Bases For Large Scale Linear Programs Exploiting Embedded Network Structure. *Computational Optimization and Applications*, **21**(1), 71-93.

Introduction

Definition 1. An **embedded network** is a structure appearing within a problem consisting of a subset of constraints and variables which define a network flow problem. The remaining constraints are referred to as **side constraints**.

- The main goal of the paper is to exploit the embedded network structure to create an advanced basis via decomposition methods.

Problem Statement

Consider the primal LP

$$\begin{aligned} \min \quad & c^T x \\ \text{s.t.} \quad & Ax = b \\ & l \leq x \leq u \end{aligned} \tag{1}$$

where $A \in \mathbb{R}^{m \times n}$, $c, x, l, u \in \mathbb{R}^n$, and $b \in \mathbb{R}^m$.

Network Decomposition

If we assume that we can detect a submatrix of network rows and columns of A is detected, and let N denote the pure network structure, we can rewrite (1) as

$$\begin{aligned}
 \min \quad & z_0 = c'^T x' + c''^T x'' \\
 \text{s.t.} \quad & Nx' = b' \\
 & Sx' + Tx'' = b'' \\
 & l' \leq x' \leq u' \\
 & l'' \leq x'' \leq u''
 \end{aligned} \tag{2}$$

where

$$\begin{aligned}
 m &= m_1 + m_2, n = n_1 + n_2, \\
 N &= [n_{ij}] \in \mathbb{R}^{m_1 \times n_1}, S = [s_{ij}] \in \mathbb{R}^{m_2 \times n_1}, T = [t_{ij}] \in \mathbb{R}^{m_2 \times n_2}, \\
 c', x', l', u' &\in \mathbb{R}^{n_1}, c'', x'', l'', u'' \in \mathbb{R}^{n_2}, b' \in \mathbb{R}^{m_1}, b'' \in \mathbb{R}^{m_2}
 \end{aligned}$$

Solving LPEN with Advanced Basis

Two things we know:

- For large-scale LPs, using an advanced basis improves sparse simplex performance
- Calculation of the initial basis is *very* important

⇒ Construct an algorithm which exploits the embedded network structure to create an advanced basis!

Decomposition Methods

Lagrangian Relaxation

- Creates a pure network flow model by adding the non-network constraints into the objective function with Lagrangian penalties
- A series of minimum cost NFPs corresponding to different values of Lagrangian multipliers are solved iteratively

Benders Decomposition

- Decomposes the LP problem into a master and a subproblem
- Cuts obtained from the subproblem are added to the master problem at each iteration

The Algorithm

- *Step 1: Preprocessing and Scaling*
- *Step 2: Network Detection*
 - Detect the network structure and decompose the problem as in 2
- *Step 3: Solve the Decomposition Problem and Create a Starting Basis*
 - Apply Lagrangean Relaxation **OR** Benders Decomposition
- *Step 4: Complete Sparse Simplex Solution*
 - Process the LPEN applying simplex method of choice using starting basis obtained in Step 3

Lagrangian Relaxation

Relaxing the side constraints $Sx' + Tx'' = b''$ yields the Lagrangian relaxation

$$\begin{aligned} \min \quad & z_{L(\lambda)} = \lambda^T b'' + (c'^T - \lambda^T S) x' + (c''^T - \lambda^T T) x'' \\ \text{s.t.} \quad & Nx' = b' \\ & l' \leq x' \leq u' \\ & l'' \leq x'' \leq u'' \\ & \lambda \in \mathbb{R}^{m_2} \text{ and unrestricted} \end{aligned} \tag{3}$$

Note that (3) is a pure NFP and can be solved efficiently using network simplex.

Lagrangian Relaxation

The dual of the the original LP (2) is given by

$$\begin{aligned}
 \max \quad & b'^T \pi + b''^T \lambda + l'^T \sigma_1 + l''^T \sigma_2 - u'^T \eta_1 - l''^T \eta_2 \\
 \text{s.t.} \quad & N^T \pi + S^T \lambda + \sigma_1 - \eta_1 = c' \\
 & T^T \lambda + \sigma_2 - \eta_2 = c'' \\
 & \sigma_1, \sigma_2, \eta_1, \eta_2 \geq 0
 \end{aligned} \tag{4}$$

- This problem is the same as the dual of (3), with different RHS values
- The Lagrangian dual problem of (2) w.r.t. the side constraints is to find multipliers λ^* that maximizes $z_{L(\lambda)}$

$$z_{L(\lambda^*)}^* = \max_{\lambda} \left\{ \min_{(x', x'')} \left(c'^T - \lambda^T S \right) x' + \left(c''^T - \lambda^T T \right) x'' \right\} \tag{5}$$

Finding the Multipliers

To determine the Lagrangean multipliers λ , we solve the LP

$$\begin{aligned} z_{L(\lambda^*)}^* &= \max w \\ \text{s.t.} \quad & w \leq f_j + \lambda^T g^j, \quad j = 1, \dots, K \end{aligned} \tag{6}$$

$$\tag{7}$$

where f_j is the objective value of (2), g^j is the subgradient of the j th basic solution

$$f_j = c'^T x'^j + c''^T x''^j, \quad g^j = b'' - Sx'^j - Tx''^j$$

and K is the number of all basic solutions of the Lagrangean problem.

Note:

- $z_{L(\lambda)}^* \leq z_0^*$ for all λ
- \exists Lagrangean multipliers λ^* such that $z_{L(\lambda^*)}^* = z_0^*$ (not true in general)

Crash Procedures

- Crash procedures are designed to create an initial basis to provide an advanced starting point
- It is well known that a starting basis with multiple structural (as opposed to logical) variables needs fewer iterations and less time to find an optimal solution
- *Triangular* crash procedures find a basis matrix which has as many structural variables as possible in such a way that the resulting basis matrix has a triangular form with a zero-free diagonal
 - The triangularity of the basis matrix ensures that a factored inverse representation of the basis with a minimum number of non-zeros can be trivially created

Network Based Crash Procedures

- For given λ the Lagrangean relaxation is a minimum cost NFP \Rightarrow feasible or optimal basis has triangular form
 - Leads to a lower triangular crash procedure considering only variables which are basic in the network optimal solution
- Construct initial basis from basic variables of the network problem and logical variables of non-network rows, apply network based crash procedure (i.e. CNET2)

Benders Motivation

The embedded NFP can be considered as a 2-stage problem in which the network and non-networks part are the 1st and 2nd stages, respectively.

⇒ Use Benders decomposition to create an advanced basis for solving the original LPEN problem

Benders Decomposition

We can split the original problem (2) into a master problem

$$\begin{aligned} \min \quad & z_M = c'^T x' \\ \text{s.t.} \quad & Nx' = b' \\ & l' \leq x' \leq u' \end{aligned} \tag{8}$$

and a subproblem

$$\begin{aligned} \min \quad & z_S = c''^T x'' \\ \text{s.t.} \quad & Tx'' = b'' - Sx'^* \\ & l'' \leq x'' \leq u'' \end{aligned} \tag{9}$$

where x'^* is an optimal solution of (8).

Subproblem Dual

The dual of the subproblem (9) is given by

$$\begin{aligned} \max \quad & \pi_1^T (b'' - Sx'^*) - \pi_2^T u'' + \pi_3^T l'' \\ \text{s.t.} \quad & \pi_1^T T \leq c'' \\ & \pi_1 \text{ free} \\ & \pi_2, \pi_3 \geq 0 \end{aligned} \tag{10}$$

The Cuts

- From duality theory, we have a feasibility cut of the form

$$\pi_1^T (b'' - Sx'^*) - \pi_2^T u'' + \pi_3^T l'' \leq 0$$

- If we let θ denote the smallest value of the upper bound of the objective function of (10), we have an optimality cut of the form

$$\theta - \pi_1^T (b'' - Sx'^*) + \pi_2^T u'' - \pi_3^T l'' \geq 0$$

Creating an Advanced Basis

Let $x^k = \left((x')^k, (x'')^k \right)$ denote the solution vector of the master and subproblem in the k th pass

- This solution may be infeasible, feasible, or feasible and optimal for the original LPEN problem
- Create the starting basis as follows
 - The variable x_i^k for the i th component appears as a basic variable in the solution of the master or subproblem \Rightarrow basic
 - If not \Rightarrow non-basic, analyze the solution values
 - * $x_i^k = l_i \Rightarrow$ non-basic at lower bound
 - * $x_i^k = u_i \Rightarrow$ non-basic at upper bound
- The basis factorization procedure INVERT uses this info to create an initial factorization of this basis as a simplex starting point for the LPEN problem

Conclusions and Future Directions

- It is shown that exploiting the embedded network structure significantly improves the performance of the simplex algorithm
 - See the paper for the computational results and comparisons among different procedures
- Refinement of decomposition methods will most likely lead to further improvement of the procedure...