# Measuring progress in branch-and-bound MILP algorithms 

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## Overview

- Reasons to measure progress of branch-and-bound
- Current measures
- Some graphical representations
- A weighted sum measure of progress
- Conclusions and future work


## Reasons to measure the progress of branch-and-bound

- How good is the best solution so far?


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- How much longer until we have a proven optimal solution?
- How likely is it that a better solution will be found, and how much better will it be?
- Should we change any algorithm strategies? (branching, node selection, cuts, ...)


## Current measures

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- Predicted tree size
- Some internal measures used for guiding the algorithm


## Optimality gap: strengths and weaknesses



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- Strength: nonincreasing


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- Strength: guarantee on quality of solution
- Strength: nonincreasing
- Weakness: may remain constant for long periods, then drop suddenly


## Number of active nodes: strengths and weaknesses



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- Weakness: not all active nodes are equal


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- The total number of nodes that will be explored is estimated early in the process
- Strength: Addresses a key question
- Weakness: Estimate is based on a common tree shape, but this tree shape depends on specific algorithm implementation and parameters


## Possible goals in analyzing b\&b progress

- Predict the time to completion


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- Predict the time to completion
- Predict the optimal objective value
- Determine when to change algorithm strategies
- Find a good measure of progress


## Information available during b\&b

- Number of active nodes


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- Number of active nodes
- For each active node:


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- Similar information for each processed node


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- So far, we have considered several instances from MIPLIB 3 that take more than 30 seconds but less than an hour


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- Usual: gap, number of active nodes
- Histogram of active node LP bounds
- Scatterplot of active node LP bounds and integer infeasibility
- Node history in scatterplot


## Histogram of active node LP bounds



- Horizontal axis is the LP bound bins
- Vertical axis is number of active nodes
- Green vertical line is the current incumbent value


## Example histogram series 1

histogram of objective values 000


## Example histogram series 1

histogram of objective values 001


## Example histogram series 1

histogram of objective values 002


## Example histogram series 1

histogram of objective values 003


## Example histogram series 1

histogram of objective values 004


## Example histogram series 1

histogram of objective values 005


## Example histogram series 1

histogram of objective values 006


## Example histogram series 1

histogram of objective values 007


## Example histogram series 1

histogram of objective values 008


## Example histogram series 1

histogram of objective values 009


## Example histogram series 1

histogram of objective values 010


## Example histogram series 1

histogram of objective values 011


## Example histogram series 1

histogram of objective values 012


## Example histogram series 1

histogram of objective values 013


## Example histogram series 1

histogram of objective values 014


## Example histogram series 2

histogram of objective values 000


## Example histogram series 2

histogram of objective values 001


## Example histogram series 2

histogram of objective values 002


## Example histogram series 2

histogram of objective values 003


## Example histogram series 2

histogram of objective values 004


## Example histogram series 2

histogram of objective values 005


## Example histogram series 2

histogram of objective values 006


## Example histogram series 2

histogram of objective values 007


## Example histogram series 2

histogram of objective values 008


## Example histogram series 2

histogram of objective values 009


## Example histogram series 2

histogram of objective values 010


## Example histogram series 2

histogram of objective values 011


## Example histogram series 2

histogram of objective values 012


## Example histogram series 2

histogram of objective values 013


## Example histogram series 2

histogram of objective values 014


## Scatterplot of active node LP bounds and integer infeasibility



- Points represent active nodes
- Vertical axis is the LP bound
- Horizontal axis is the sum of integer infeasibilities
- Green horizontal line is the current incumbent value


## Example scatterplot series

scatterplot 000


## Example scatterplot series

scatterplot 001


## Example scatterplot series

scatterplot 002


## Example scatterplot series

scatterplot 003


## Example scatterplot series

scatterplot 004


## Example scatterplot series

scatterplot 005


## Example scatterplot series

scatterplot 006


## Example scatterplot series

scatterplot 007


## Example scatterplot series

scatterplot 008


## Example scatterplot series

scatterplot 009


## Example scatterplot series

scatterplot 010


## Example scatterplot series

scatterplot 011


## Example scatterplot series

scatterplot 012


## Example scatterplot series

scatterplot 013


## Example scatterplot series

scatterplot 014


## History of active nodes



- Shows the ancestors of the node


## Example histories of active nodes

path of incumbents 000


## Example histories of active nodes

path of incumbents 001


## Example histories of active nodes

path of incumbents 002


## Example histories of active nodes

path of incumbents 003


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- The user would like a good measure of progress
- Gap and number of active nodes don't work well
- Histograms give good information, but we want a single value
- One idea: sum of gaps
- But this fluctuates a great deal (with the number of active nodes)
- Another idea: average gap


## Measure of progress: Average gap doesn't work well

histogram of objective values 001


## Measure of progress: Average gap doesn't work well

histogram of objective values 002


## Measure of progress: Average gap doesn't work well

histogram of objective values 003


## Measure of progress: Average gap doesn't work well

histogram of objective values 004


## Measure of progress: Average gap doesn't work well

histogram of objective values 005


## Measure of progress: Average gap doesn't work well

histogram of objective values 006


## Measure of progress: Weighted sum of active node gaps

- Current measure: Weight each node based on depth. Let $A$ be the set of active nodes, $g_{i}$ be the gap for node $i$, and depth $h_{i}$ be the depth of node $i$ :

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\sum_{i \in A} \frac{g_{i}}{2^{d_{i}}}
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- Valuable properties:
- Sum of weights of children equals parent's weight
- Weights are constant
- Therefore: Monotonic decreasing (as long as lp bounds of parent and child differ)


## Example graphics




## Example graphics



I152lav, CBC default


## Example graphics


stein45, CBC default


## Example graphics


misc07, CBC default


## Example graphics


bell3a, GLPK intopt no cuts


## Example graphics


bell5, GLPK intopt no cuts


## Example graphics


bell3a, GLPK standard, best bound


## Example graphics


misc07, GLPK standard


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bell3a, CBC default


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stein45, GLPK intopt with cuts


## Strengths and weaknesses

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- monotonic decreasing whenever child LP bound differs from parent
- generally smoother measure of progress
- appears robust to different solvers and options
- Weakness: still drops significantly when new incumbents found


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- However, the predictions do not appear to be consistently accurate, especially for big drops


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- Valuable to represent data visually when considering summary measures
- We should explore more data mining/machine learning applied to MIPs
- Value of open-source codes: proven useful on real-world problems and allow full and easy access to information available during the algorithm


## Current and Future work

- Finish examining measures of progress and publish these ideas, including code to generate graphics


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- Can other information be extracted: recommended node selection strategy or cuts?

